

Addition Automata and Attractors of Digit Systems Corresponding to Expanding Rational Matrices

Anjelo Gabriel R. Cruz¹
Manuel Joseph C. Loquias¹
Jörg M. Thuswaldner²

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¹ University of the Philippines, Diliman

² Montanuniversität Leoben

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- (A, \mathcal{D}) has the finiteness property: every $x \in \mathbb{Z}^n$ has an expansion of the form

$$x = d_0 + Ad_1 + \cdots + A^{k-1}d_{k-1} \quad (1)$$

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- 1 $\mathbb{Z}^n[A]$ = smallest nontrivial A -invariant \mathbb{Z} -submodule of \mathbb{Q}^n containing \mathbb{Z}^n
- 2 $\mathbb{Z}^n[A]/A\mathbb{Z}^n[A]$ is a finite abelian group (Jankauskas and Thuswaldner, 2022)

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Remark. $\forall x \in \mathbb{Z}^n[A], \forall k \in \mathbb{Z}^+,$

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where $d_i = \mathbf{d}(\Phi^i(x))$ for all $i < k$

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③ $n = 1 \implies A = \frac{p}{q} \in \mathbb{Q}$ in lowest terms with $|p| > |q|$
(Akiyama, Frougny, and Sakarovitch 2008)

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where $\varepsilon \in \{1, -1\}$, $p(x) = x^2 + \alpha x + \beta \in \mathbb{Z}[x]$ is irreducible, and $\mathcal{D} = \{(0, 0)^\top, (1, 0)^\top, \dots, (|\beta| - 1, 0)^\top\}$

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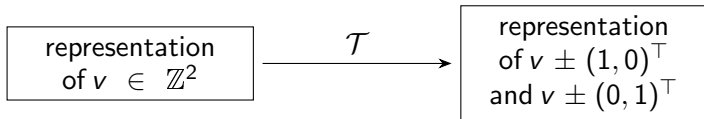
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Remark.

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\therefore every $x \in \mathbb{Z}^n[A]$ can be written uniquely in the form

$$\begin{aligned} x &= \sum_{j=0}^{k-1} A^j d_j + \sum_{\ell=0}^{Nr-1} A^{k+\ell} b_{\ell \bmod r} + A^{Nr+k} p \quad (\forall N \in \mathbb{Z}^+) \\ &=: (\sigma)_A, \end{aligned}$$

where $\sigma = (b_{r-1} \cdots b_1 b_0)^\infty d_{k-1} \cdots d_1 d_0$ (**A-adic representation** of x)

Remark.

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- ❷ $x = (0^\infty d_{k-1} \cdots d_1 d_0)_A$
 $\implies x$ can be identified with the finite word $d_{k-1} \cdots d_1 d_0$

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Proposition

$\mathcal{D} = \{(0, 0)^\top, (1, 0)^\top, \dots, (|b| - 1, 0)^\top\}$ is a complete residue system of $\mathbb{Z}^2[A]/A\mathbb{Z}^2[A]$.

Given:

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Define the following operations on σ :

$$(\sigma^{\pm P})_A = (\sigma)_A \pm (1, 0)^\top$$

$$(\sigma^{\pm Q})_A = (\sigma)_A \pm A(c, 0)^\top \pm (a - c, 0)^\top$$

$$(\sigma^{\pm R})_A = (\sigma)_A \mp A(c, 0)^\top \mp (a, 0)^\top$$

$$(\sigma^{\pm S})_A = (\sigma)_A \pm (c, 0)^\top$$

$$(\sigma^{\pm T})_A = (\sigma)_A \pm A(c, 0)^\top \pm (a + c, 0)^\top$$

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$$\therefore \pm U \xrightarrow{\nu|\nu} \pm P$$

The case $0 < \alpha \leq \beta - 1$

Given:

- expanding matrix $A = \begin{bmatrix} 0 & -\beta \\ 1 & -\alpha \end{bmatrix} \in \mathbb{Q}^{2 \times 2}$ with irreducible characteristic polynomial $p(x) = x^2 + \alpha x + \beta \in \mathbb{Q}[x]$
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Theorem (C.-Loquias-Thuswaldner, preprint, 2025)

Given: $0 < \alpha \leq \beta - 1$

Then: $\mathcal{A}_\Phi = \{0\}$

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Notation.

- ① $(\nu, 0)^\top \in \mathcal{D}$ will be identified with ν
- ② $\overline{f(\nu)}$: $f(\nu) \in \mathcal{D}$ and is taken to be a single digit in the expansion

Transition relations for $\nu \in \mathcal{D}$:

$$(\sigma\nu)^P = \begin{cases} \sigma \overline{\nu + 1}, & 0 \leq \nu < b - 1 \\ \sigma^R 0, & \nu = b - 1 \end{cases}$$

$$(\sigma\nu)^{-P} = \begin{cases} \sigma \overline{\nu - 1}, & 1 \leq \nu < b \\ \sigma^{-R} \overline{b - 1}, & \nu = 0 \end{cases}$$

$$(\sigma\nu)^{\pm Q} = \begin{cases} \sigma^{\pm T} \overline{\nu \pm a \mp c \pm b}, & 0 \leq \nu \pm a \mp c \pm b < b \\ \sigma^{\pm S} \overline{\nu \pm a \mp c}, & 0 \leq \nu \pm a \mp c < b \\ \sigma^{\mp Q} \overline{\nu \pm a \mp c \mp b}, & 0 \leq \nu \pm a \mp c \mp b < b \end{cases}$$

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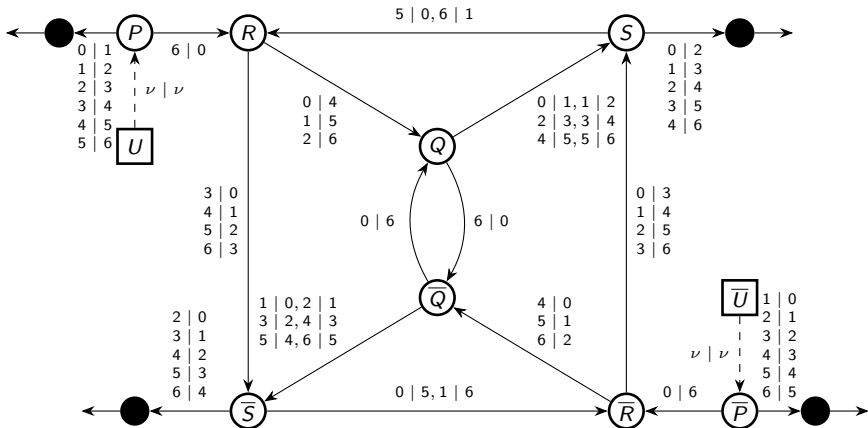
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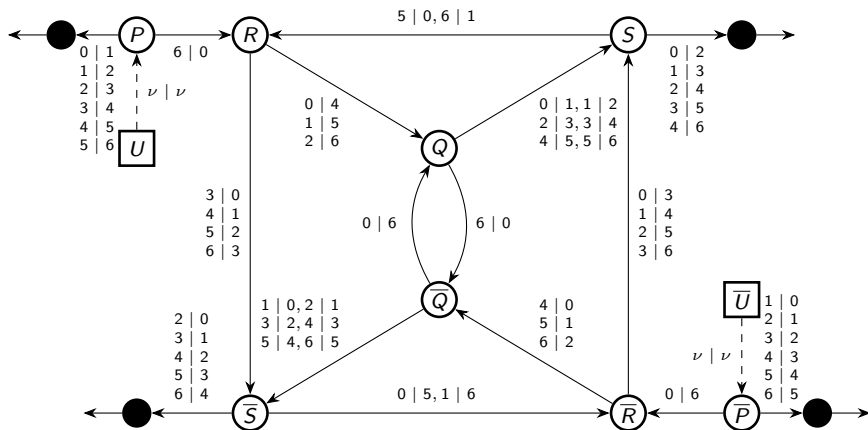
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Transition diagram for $\alpha = \frac{3}{2}$ and $\beta = \frac{7}{2}$



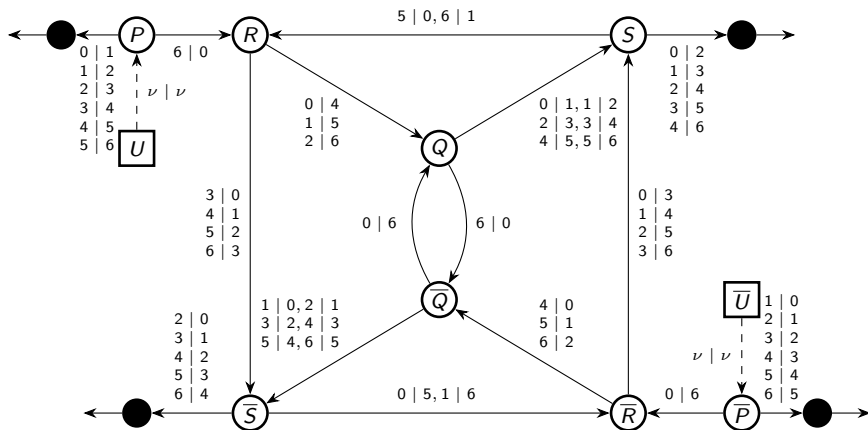
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Remark. Any walk in \mathcal{T} with label of the form $0 \mid d$ always leads to \bullet .

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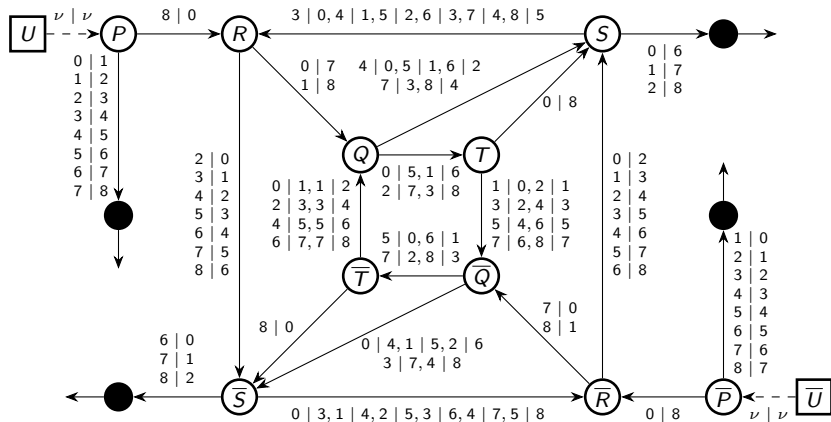
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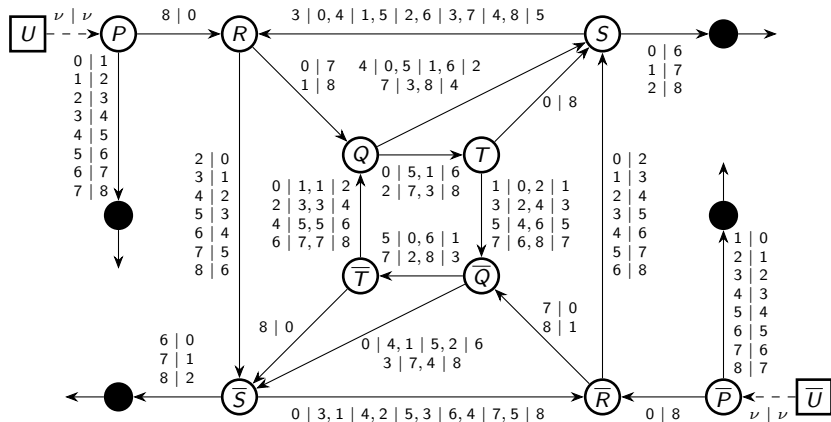
Transition diagram for $\alpha = \frac{1}{3}$ and $\beta = \frac{3}{2}$ (note: $0 < \alpha < 1$):



e.g.: $v = (17, 0)^T$ with representation $\sigma = 68578$

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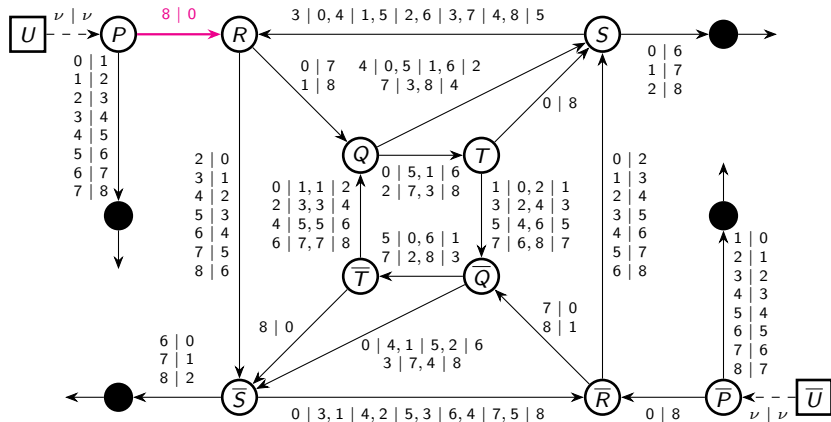


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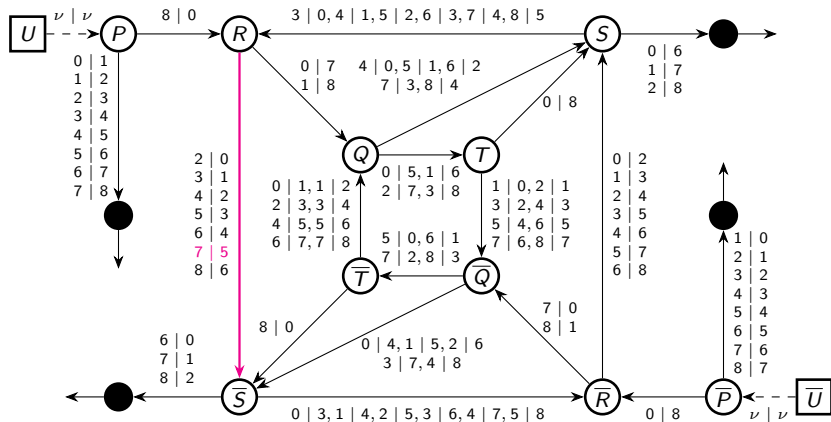
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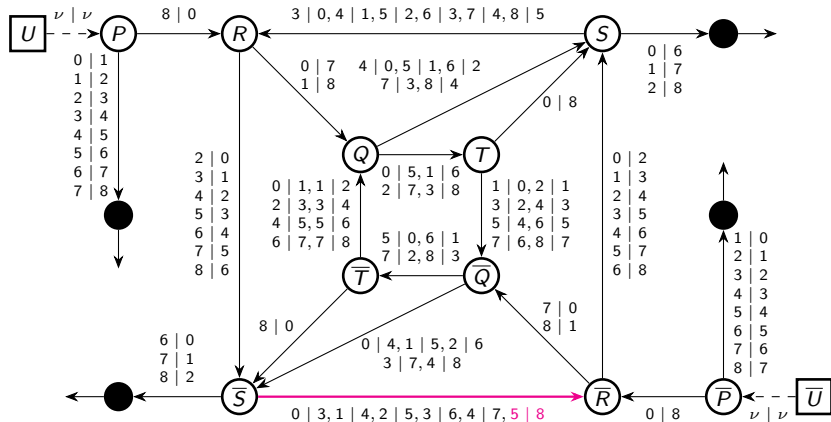


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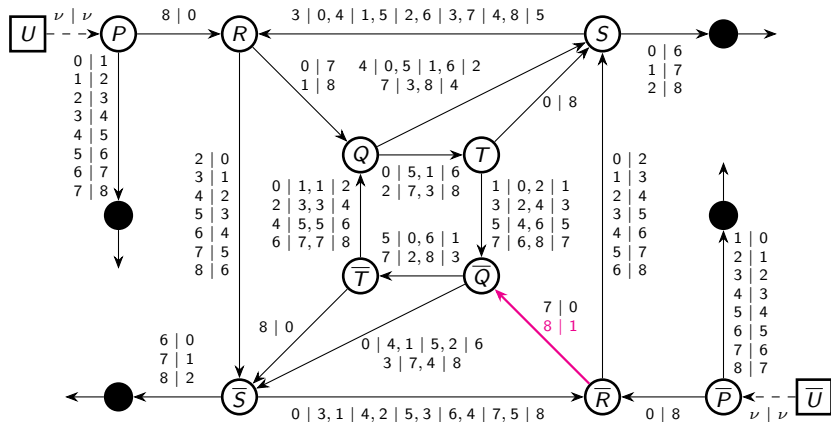


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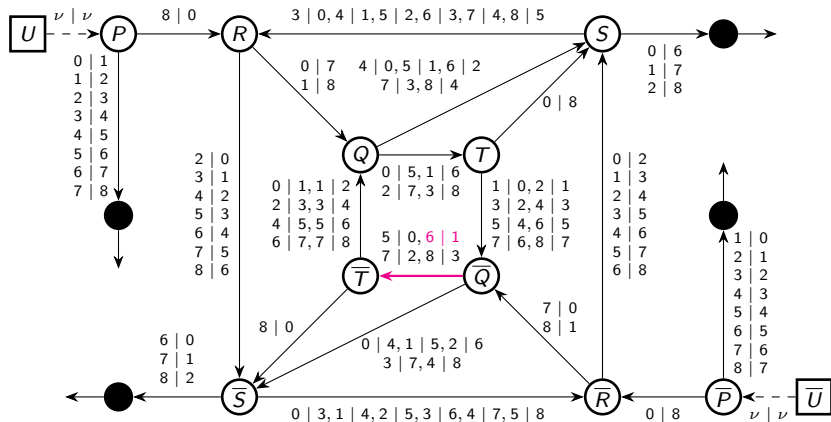


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\Rightarrow representation of $(18, 0)^T = (\sigma^P)_A$: $\sigma^P = 1850$

The case $0 < \alpha \leq \beta - 1$

Transition diagram for $\alpha = \frac{1}{3}$ and $\beta = \frac{3}{2}$ (note: $0 < \alpha < 1$):

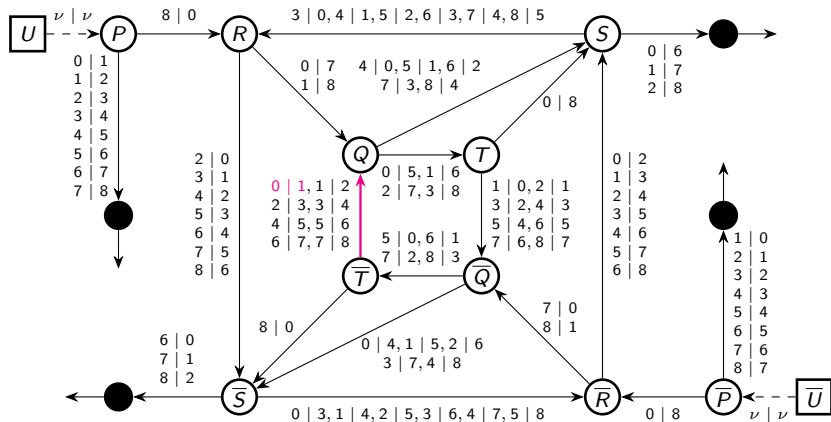


e.g.: $v = (17, 0)^T$ with representation $\sigma = 68578$

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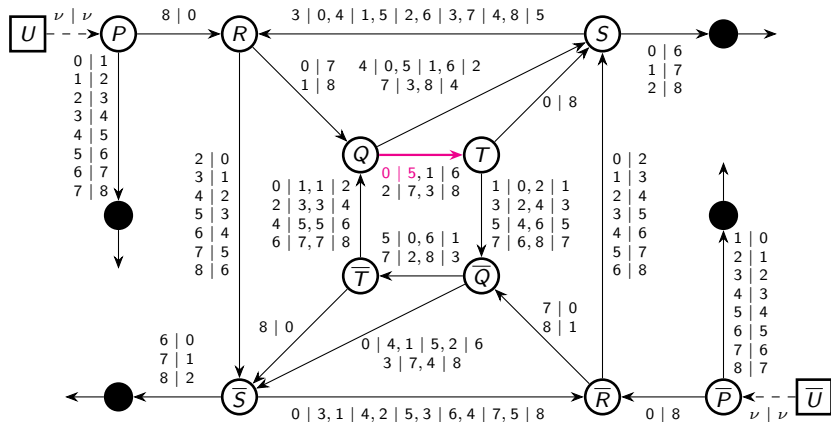


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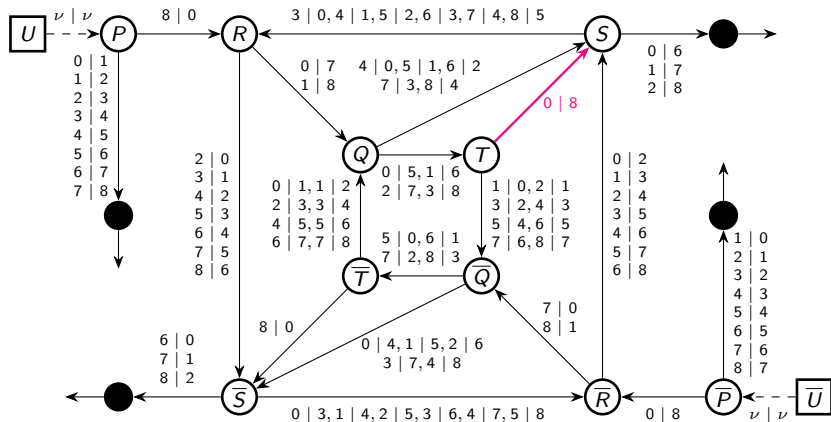


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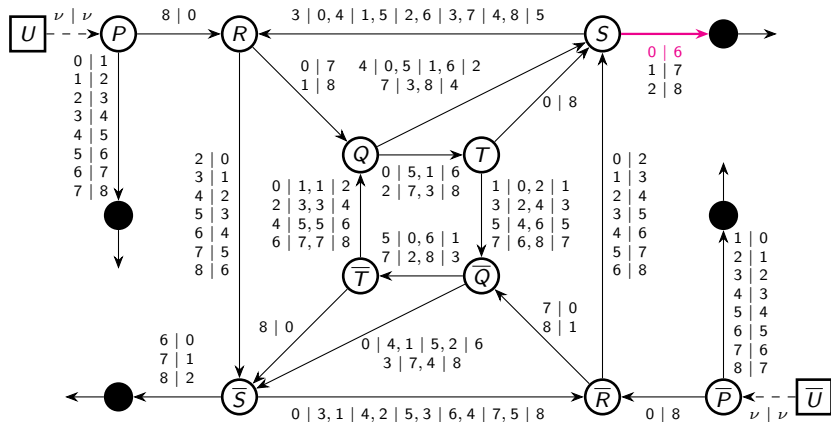


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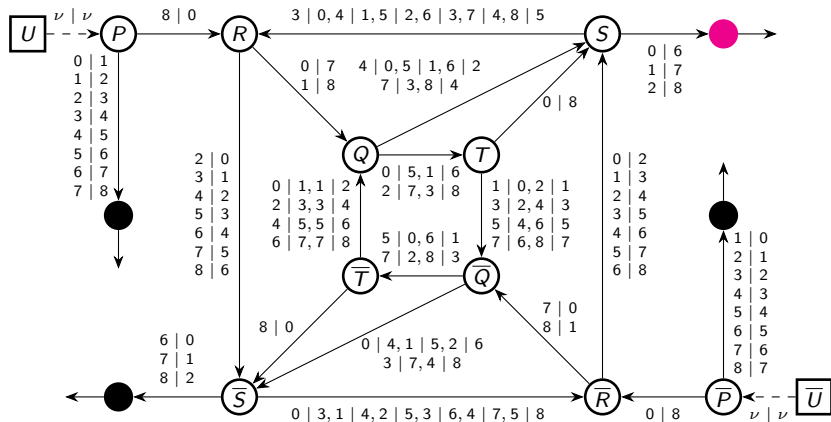


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Feed $(0, 0)^\top$ (with representation $\sigma = 0^\infty$) into \mathcal{T} at $\pm P$ and $\pm U$.

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The case $0 < -\alpha < \beta - 1$

Theorem (C.-Loquias-Thuswaldner, preprint, 2025)

Given: $0 < -\alpha < \beta - 1$

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Theorem (C.-Loquias-Thuswaldner, preprint, 2025)

Given: $0 < -\alpha < \beta - 1$

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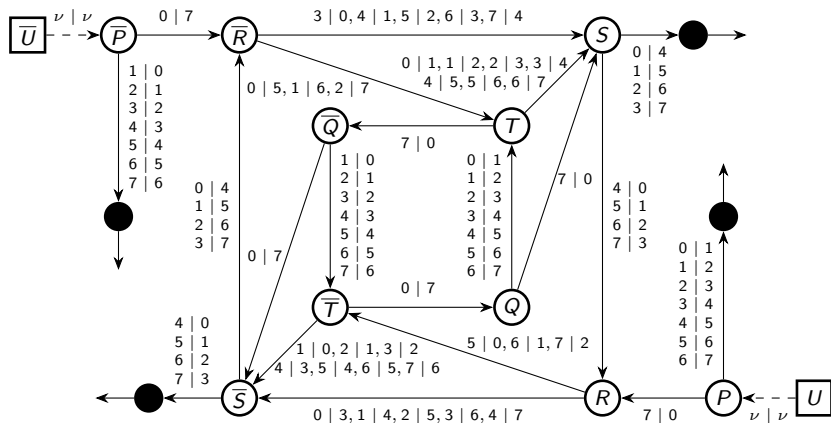
Then:

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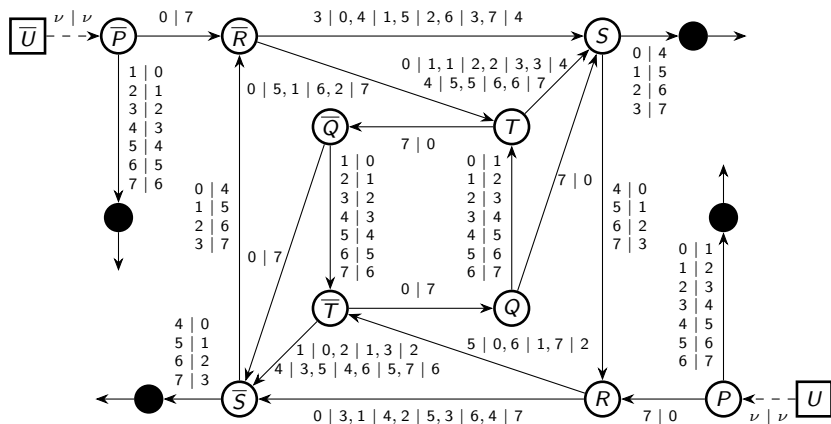
The case $0 < -\alpha < \beta - 1$

Transition diagram for $\alpha = -\frac{3}{4}$ and $\beta = 2$ (note: $\alpha \geq -1$):



The case $0 < -\alpha < \beta - 1$

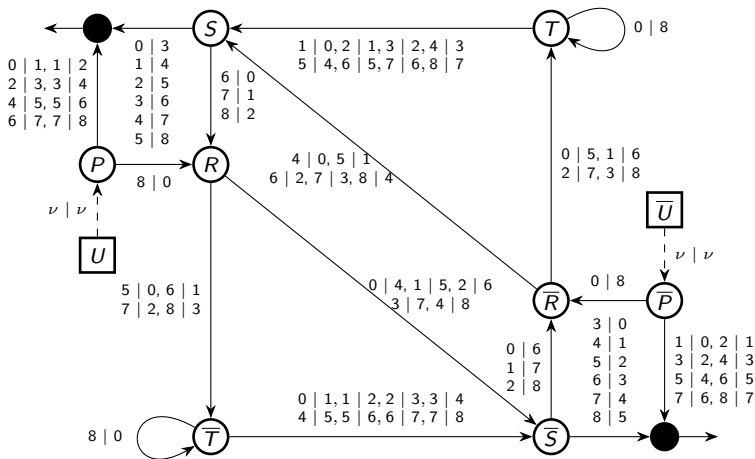
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Remark. Any walk in \mathcal{T} with only 0 as input always leads to \bullet .

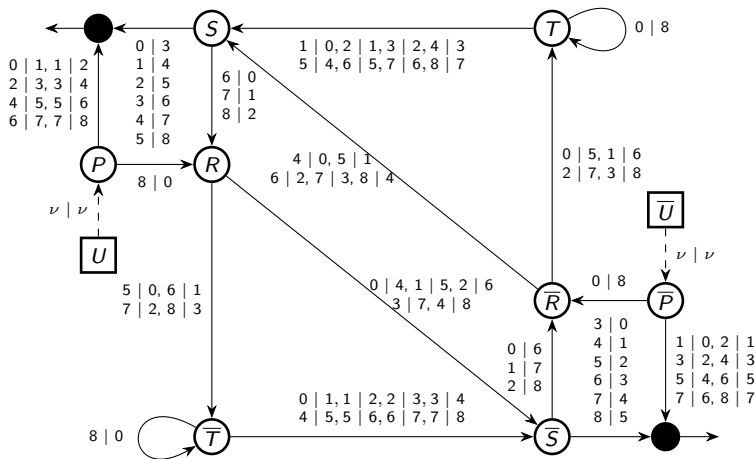
The case $0 < -\alpha < \beta - 1$

Transition diagram for $\alpha = -\frac{4}{3}$ and $\beta = 3$ (note: $\alpha < -1$):



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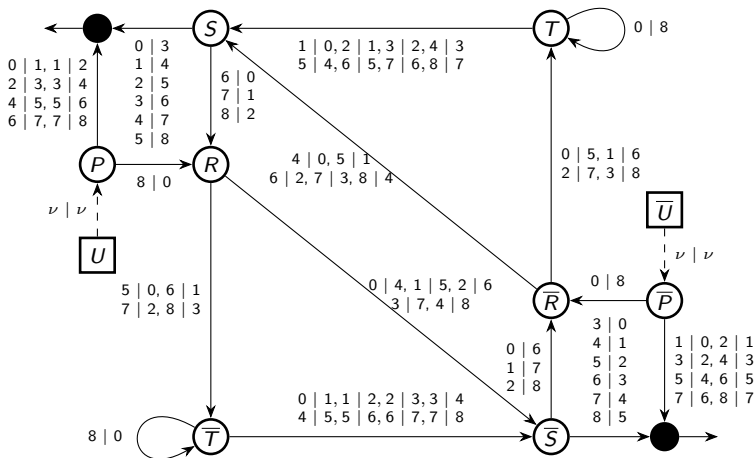
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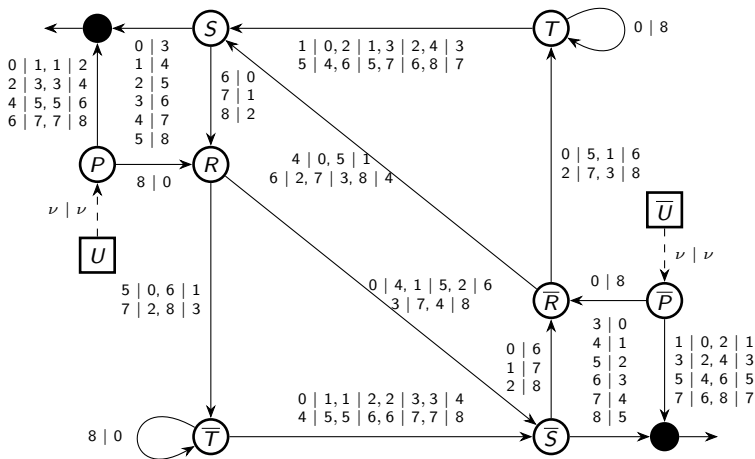
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e.g.: $\nu = (8, 2)^\top$ with representation $\sigma = 28$

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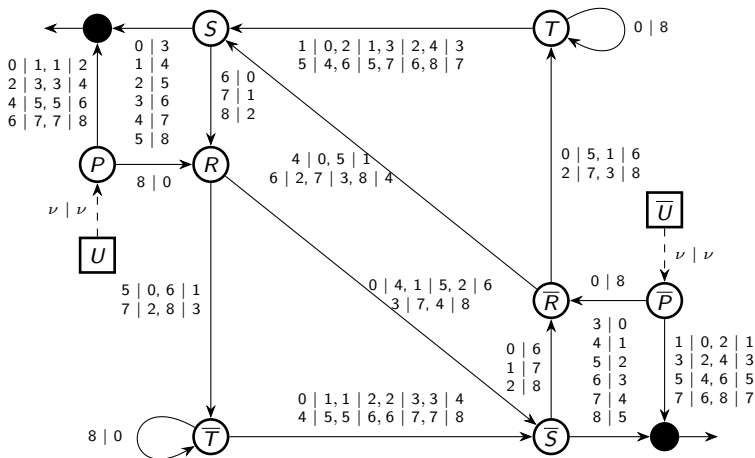


e.g.: $\nu = (8, 2)^\top$ with representation $\sigma = 28$

\Rightarrow representation of $(9, 2)^\top = (\sigma^P)_A$: $\sigma^P = 8^\infty 5660$

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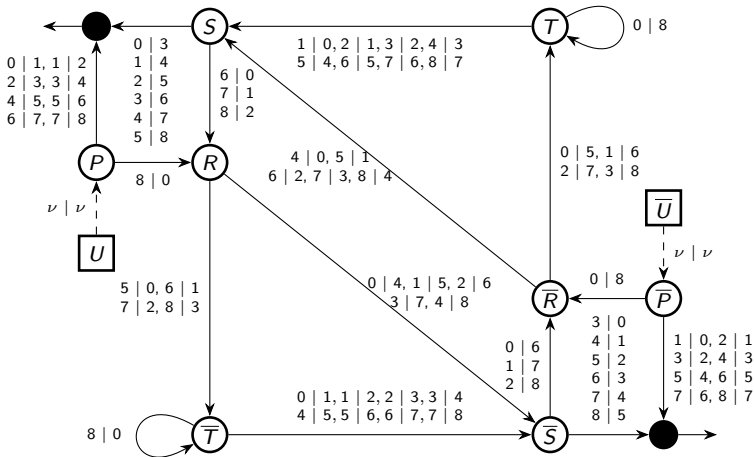


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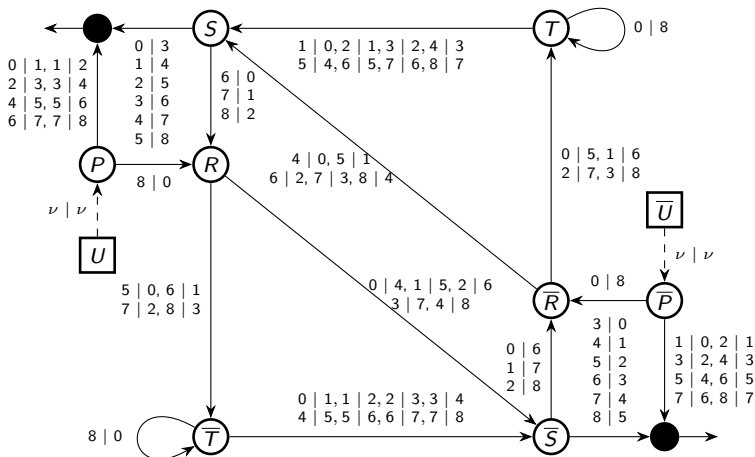
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$$\implies p := (8^\infty)_A \in \mathcal{A}_\Phi$$

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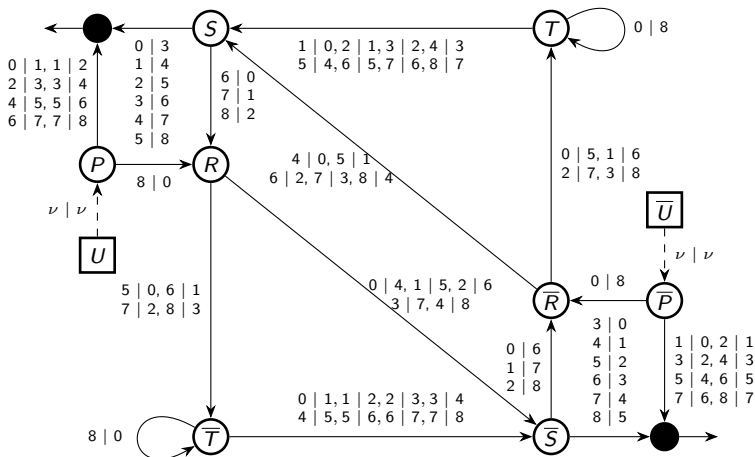


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Theorem (C.-Loquias-Thuswaldner, preprint, 2025)

Given: $0 < -\alpha < -\beta - 1$

Then: \mathcal{A}_Φ **contains** $\{(0, 0)^\top, -(a, c)^\top, -(c, 0)^\top\}$

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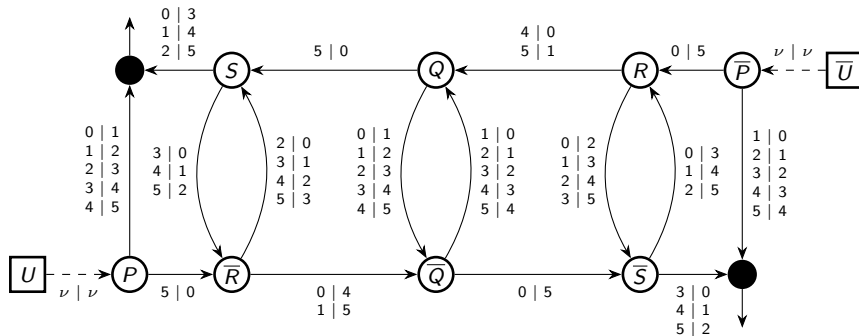
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Remark. \mathcal{A}_Φ may contain other points depending on α and β .

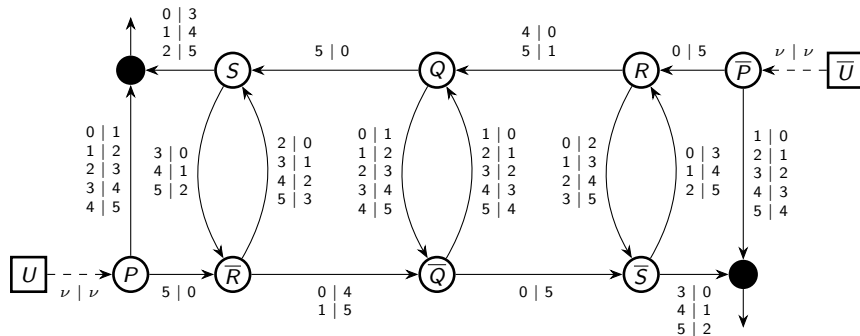
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Transition diagram for $\alpha = -\frac{2}{3}$ and $\beta = -2$:



The case $0 < -\alpha < -\beta - 1$

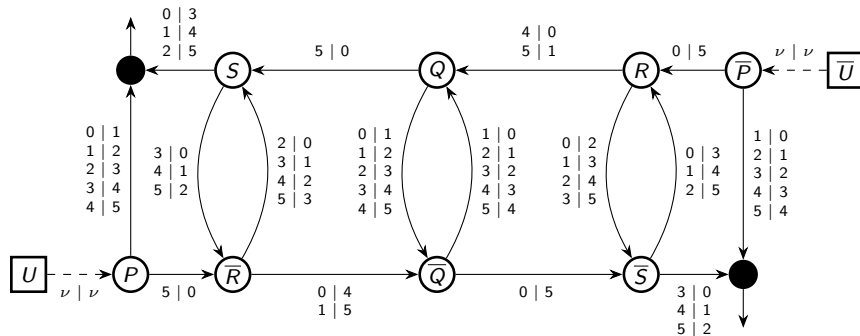
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Remark. Any walk in \mathcal{T} with only 0 as input leads to one of the following:

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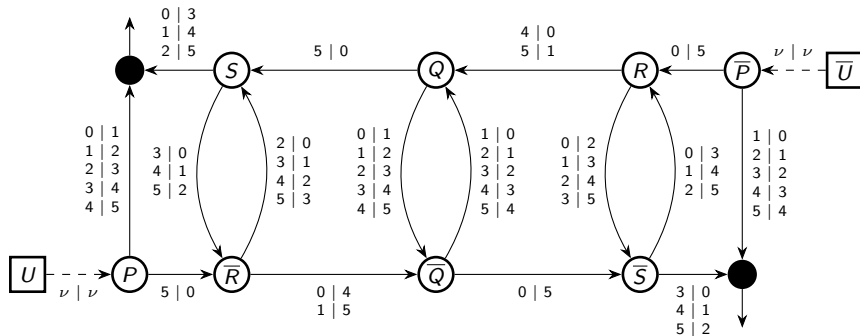


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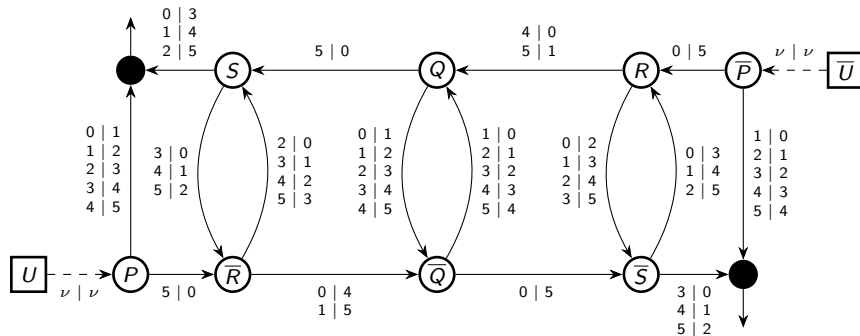


Remark. Any walk in \mathcal{T} with only 0 as input leads to one of the following:

- 1 the terminal state \bullet
- 2 the loop $R \rightarrow -S \rightarrow R$

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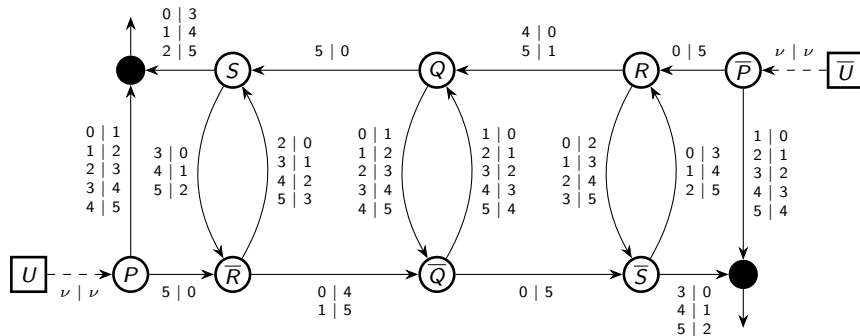


e.g.: $v = (7, -3)^\top$ with representation $\sigma = (32)^\infty 05$

\implies representation of $v + (1, 0)^\top = (\sigma^P)_A$ is $\sigma^P = (41)^\infty 40$

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Transition diagram for $\alpha = -\frac{2}{3}$ and $\beta = -2$:



e.g.: $v = (7, -3)^\top$ with representation $\sigma = (32)^\infty 05$

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and $p = (41)^\infty = (1, 6)^\top \in \mathcal{A}_\Phi$

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Given: $0 < -\alpha < \beta - 1$

$$\gamma = -a - b - c$$

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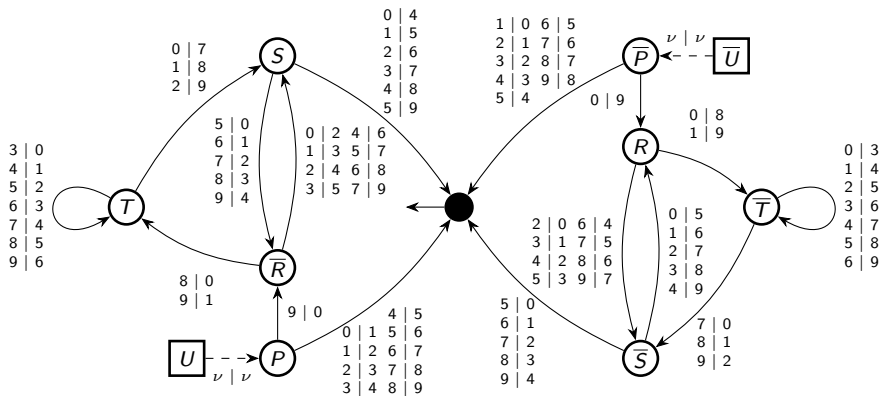
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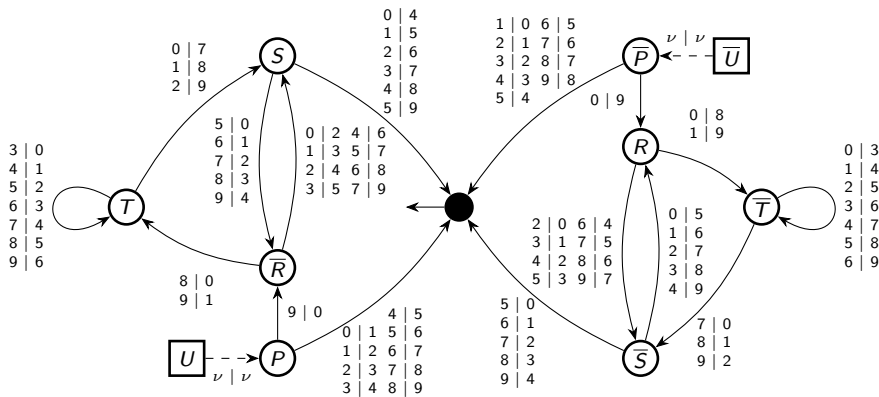
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Transition diagram for when $\alpha = \frac{2}{5}$ and $\beta = -2$:



The case $0 < \alpha < -\beta - 1$

Transition diagram for when $\alpha = \frac{2}{5}$ and $\beta = -2$:



Remark. Any walk in \mathcal{T} with only 0 as input may lead to either the terminal state \bullet or the self-loop at $-T$

Theorem (C.-Loquias-Thuswaldner, preprint, 2025)

① $\alpha > 1$ and $\beta - 1 \leq \alpha \leq \beta \implies \mathcal{A}_\Phi = \{0\}$

② $0 < \alpha \leq 1$ and $0 < \beta - 1 < \alpha < \beta$

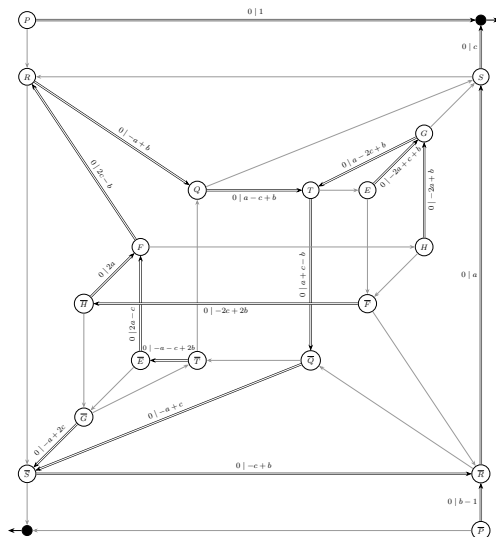
$\implies \mathcal{A}_\Phi = \{0\}$ *provided one of the following holds:*

① $\beta > 2\alpha$ and $\alpha + \beta > 2$

② $\beta \leq 2\alpha$ and $\alpha + 2 < 2\beta$

Additional Cases

Transition diagram for Case 2.1:



Transition diagram for Case 2.2: